

# A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW



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**ABSTRACT:** Networks play an important role in electrical and electronic engineering. It depends on what area of electrical and electronic engineering, for example, there is a lot more abstract mathematics in communication theory and signal processing and networking, etc. Networks involve nodes communicating with each other. Graph theory has found considerable use in this area. In this paper, we introduce some new Networks such as Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW, and Change Network Symmetric Digraph- PW. Moreover, several theorems and results of these networks have been studied.

**KEYWORDS:** Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW, Change Network Symmetric Digraph- PW.

## INTRODUCTION

For more detail on Graphs, Digraphs, and network Digraphs one. The networks (Graph- *PW* and symmetric Digraph- *PW*) and network changes (Graph-*PW* and symmetric Digraph-*PW*) can be the source of many algorithms of practical importance. It is flexibly adaptable to suit the needs of the application, so it can be used in areas such as syntactic analysis, fault detection, and diagnosis in computer, therefore the graphical of this representation of the object and the binary relation on them is a convenient form of expression. In this paper, we give some definitions and results of network Graphs-*PW* and symmetric Digraphs-*PW* and networks changes (Graph-*PW* and symmetric Digraph-*PW*) [1,2], we refer to the interested reader to [3-30].

### 1. Networks (Graph-*PW* And Symmetric Digraph-*PW*):

**Definition (1.1):** A **Graph-*PW*** is a triple  $G_{PW} = (V, E^*, \Phi)$  consists of a non-empty set  $V = \{v_1, v_2, \dots, v_n\}$  of objects called vertices, points, nodes, or just dots, together with undirected pairs set of vertices

$$E^* = \left\{ \begin{array}{l} e_{ij} = v_i v_j = v_j v_i : i \neq j \\ \text{or } i = j, n(e_{ij}) = v_i * v_j = \rho\omega \end{array} \right\}$$

is called edges, arcs or lines, satisfy  $PW = \max(n(e_{ij}) = v_i * v_j = \rho\omega)$ , where  $P$  (non-negative integer) is the maximum numbers of  $\rho$ -edges or loops between any pairs of vertices and  $W$  is the uniform weighted with  $PW = \max \rho\omega$  depended on  $\omega$  uniform sign weighted in  $E^*$ , and incident function  $\Phi$  from  $E^*$  to the set  $P_2(V)$  of all 2-elements or parts subsets of  $V$ , that is,  $\Phi: E^* \mapsto P_2(V)$ . The **adjacency function matrix  $Am(G_{PW})$**  define as

$$Am(G_{PW}) = \begin{cases} n(e_{ij}) \text{ if } \rho > 0 \wedge \omega > 0 \wedge v_i \text{ joined to } v_j; \\ 0 \text{ or } \infty \text{ if } \rho\omega = 0 & ; \\ -n(e_{ij}) \text{ if } \rho > 0 \wedge \omega < 0 \wedge v_i \text{ joined to } v_j. \end{cases}$$

clear that it is if  $P = W = 1$ , then the Graph-*PW* is natural **Graph-1**, and if  $P > 0, W = 1$  the Graph-*PW* is **mullet Graph- $P$** , and if  $P = 1, W \in \mathbb{R}^+$  the Graph-*PW* is **weighted Graph- $W$**  and  $\Phi$  in Graph-*PW* satisfy  $\forall e_{ij} \in E^* \exists \{v_i, v_j\} \in P_2(V) \ni \Phi(e_{ij}) = \{v_i, v_j\}$ ,  $\Phi$  is **one -to-one** if and only if  $\Phi(e_{ij}) = \Phi(e_{rs}) \Rightarrow e_{ij} = e_{rs}$  or  $e_{ij} \neq e_{rs} \Rightarrow \Phi(e_{ij}) \neq \Phi(e_{rs})$ , **onto** if  $\forall \{v_i, v_j\} \in$

## A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW

$P_2(V) \exists e_{ij} \in E^* \ni \Phi(\{v_i, v_j\}) = e_{ij}$  or  $\Phi(P_2(V)) = E^*$ , and if  $\Phi$  one-to-one and onto is called **corresponding** then there is an inverse  $\Phi^{-1}$  of  $\Phi$ .

The **operation** \* is modules, plus, product, max, min or ... and the **examples** are only define the operation \* by these operations. In the following definitions and theorems  $PW$ , incident function and adjacency function matrix define as well as.

**Definition (1.2):** In the **network Graph- PW**, the weight  $\omega$  is called **flow function** for each edges of  $\rho \in n(e_{ij})$  with  $\omega_i \leq \omega_{i+1}$  and  $\rho = \sum_{i=1}^{\rho} \omega_i$  between any adjacent two pair vertices is called **capacity constraint** with  $\omega \leq \rho = \sum_{i=1}^{\rho} \omega_i$  and the total flow function for any vertex is  $\omega(e) = \sum_{i=0} \omega_{i+1}$ , and the total capacity constraint for any vertex is  $\rho(e) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j$  with  $\omega(e) \leq \rho(e)$ . If  $\omega(e) = \sum_{i=0} \omega_{i+1} = PW$  then  $PW$  is called **value flow**, therefore, for any intermediate vertex if the total flow is  $PW$  of these vertices called **flow conservation**.

By these arguments the name **Graph- PW** and **network Graph- PW** are the same.

**Definition (1.3):** The **complement** of  $G_{PW}$  is  $\bar{G}_{PW} = (V, \bar{E}^*, \bar{\Phi})$  with  $V(G_{PW}) = V(\bar{G}_{PW})$ , and  $\bar{E}^* = \{e_{ij} = v_i v_j \wedge \bar{n}(e_{ij}) = PW - n(e_{ij}) \forall v_i, v_j \in V(G_{PW})\}$ , and  $\bar{\Phi}: \bar{E}^* \rightarrow P_2(V)$ .

Therefore  $G_{PW} \cup \bar{G}_{PW} = K_{|V|} - PW$  is called **complete Graph- PW**.

**Definition (1.4):** The **regular** of  $G_{PW}$  is  $R_{PW} = (V, E^*, \Phi)$  with

$$E^* = \{e_{ij} = v_i v_j : \deg v_i = \deg v_j \wedge n(e_{ij}) = v_i * v_j = \rho\omega = PW \forall v_i, v_j \in V(G_{PW})\}.$$

$G_{PW}$  is **semi-regular** if one or two vertices have the equal degree different of all vertices, that is,

$$(\deg v_1 = \deg v_2) \neq (\deg v_i = \deg v_j) \forall i, j \wedge n(e_{ij}) = v_i * v_j = PW$$

**Theorem (1.1):** If a Graph- PW  $G_{PW} = (V, E^*, \Phi)$  or  $\bar{G}_{PW}$  has an isolated vertex, then  $G_{PW} - v$  or  $G_{PW} \cup \bar{G}_{PW}$  has not an isolated vertex.

**Proof:** Since  $G_{PW} = (V, E^*, \Phi)$  define on operations, then the operation module has only isolated vertex, and the operation of product has an isolated vertex, if one of vertex is zero, and the other operations  $\bar{G}_{PW}$  has an isolated vertex, therefore  $G_{PW} - v$  or  $G_{PW} \cup \bar{G}_{PW}$  has not an isolated vertex.

Now, by these theorem  $G_{PW}$ ,  $\bar{G}_{PW}$  and  $G_{PW} \cup \bar{G}_{PW}$  well define are connected and we can be taken these networks.

**Definition (1.5):** In the value flow of the networks, if  $\sum_{i=0} \omega_{i+1} > PW$ , then there is  $\sum_{j=0} \omega_{j+1}$  such that  $\sum_{i=0} \omega_{i+1} - \sum_{j=0} \omega_{j+1} = PW$ , and  $PW \leq \rho(e) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j$ .

**Definition (1.6):** A **sub-Graphs- PW** of a Graph- PW is  $H_{pw}$  or  $\hat{G}_{PW} = (V, E^*, \Phi) \subseteq G_{PW}$  with  $V(H_{pw}) \subseteq V(G_{PW})$  or  $V(\hat{G}_{PW}) = V(G_{PW})$ ,  $E^*(H_{pw})$  or

$$E^*(\hat{G}_{PW}) \subseteq E^*(G_{PW}), \text{ and } \Phi(H_{pw}) \text{ or } \Phi(\hat{G}_{PW}) \subseteq \Phi(G_{PW}).$$

Now, let  $H_{pw} = (V(H), E^*(H), \Phi(H))$ ,  $Y_{pw} = (V(Y), E^*(Y), \Phi(Y)) \subseteq G_{PW} = (V, E^*, \Phi)$  with  $V(Y) = V(G_{PW}) - V(H)$ ,  $E^*(Y) = \{e_{ij} = v_i v_j : n(e_{ij}) = v_i * v_j = \rho\omega \notin E^*(H)\}$ . Then  $H_{pw} \cup Y_{pw} \subseteq G_{PW}$ , and  $H_{pw} + Y_{pw} = G_{PW}$ , where  $V(H) \cup V(Y) = V(G_{PW})$  and  $E^*(H) \cup E^*(Y) \cup \{e_{ij} = v_i v_j : v_i \in V(H), v_j \in V(Y) \forall v_i, v_j, n(e_{ij}) = v_i * v_j = \rho\omega\}$ .

Let  $H_{pw} = (V(H), E^*(H), \Phi(H))$ ,  $Y_{pw} = (V(Y), E^*(Y), \Phi(Y)) \subseteq \hat{G}_{PW} = (V, E^*, \Phi) \subseteq G_{PW}$  with  $V(Y) = V(\hat{G}_{PW}) - V(H)$ ,  $E^*(Y) = \{e_{ij} = v_i v_j : n(e_{ij}) = v_i * v_j = \rho\omega \notin E^*(H)\}$ . Then  $H_{pw} \cup Y_{pw} \subseteq \hat{G}_{PW}$ , and

$$H_{pw} \cup Y_{pw} = \hat{G}_{PW} \text{ iff}$$

$$E^*(H) \cup E^*(Y) \cup \{e_{ij} = v_i v_j : v_i \in V(H), v_j \in V(Y) \wedge n(e_{ij}) = v_i * v_j = \rho\omega \in E^*(\hat{G}_{PW})\}.$$

Now, we can give the following results on the network Graph- PW on flow function, value flow and capacity constraint.

**Theorem (1.2):** If  $H_{pw} = (V(H_{pw}), E^*(H_{pw}), \Phi(H_{pw})) \subseteq G_{PW} = (V, E^*, \Phi)$  of network Graph- PW, with  $pw = PW$ , then  $H_{pw} + G_{PW} - V(H_{pw}) = G_{PW}$ .

**Proof:** By definition (1.6) the result follows.

## A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW

**Theorem (1.3):** Let  $\omega(e) = \sum_{i=0} \omega_{i+1}$ ,  $\rho(e) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j > PW$  be total flow function and capacity constraint of networks Graph- PW. Then there is

$\sum_{r=0} \omega_{i+1}(e), \rho(e) = \sum_{r=1} (\sum_{i=1}^{\rho} \omega_i)_r \in H_{pw}$  or  $G_{PW} - V(H_{pw}) \subseteq G_{PW} = (V, E^*, \Phi)$  such that

- I.  $\omega(e) - \sum_{r=0} \omega_{r+1} = PW$ , and
- II.  $\rho(e) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j - \rho(e) = \sum_{r=1} (\sum_{i=1}^{\rho} \omega_i)_r = PW$ .

**Proof:** I. Since  $\sum_{r=0} \omega_{i+1} \in H_{pw}$  or  $G_{PW} - V(H_{pw})$ , then  $\sum_{r=0} \omega_{r+1} \in G_{PW}$ , and by definition (1.5) the result follows, and

- II. Since  $\rho(e) = \sum_{r=1} (\sum_{i=1}^{\rho} \omega_i)_r \in H_{pw}$  or  $G_{PW} - V(H_{pw})$ , then  $\rho(e) = \sum_{r=1} (\sum_{i=1}^{\rho} \omega_i)_r \in G_{PW}$ , and the result follows.

**Theorem (1.4):** Let  $\omega(e) = \sum_{i=0} \omega_{i+1}$ ,  $\rho(e) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j$  be total flow function and capacity constraint of networks Graph- PW have value flow and let,  $\omega_j(e), \rho_j(e)$  be total flow function and capacity constraint of sub-graph- pw from  $H_{pw}$  to  $G_{PW} - V(H_{pw})$ . Then  $PW = \omega_j(e) - \omega_r(e), PW = \rho_j(e) - \rho_r(e)$ , where  $\omega_r(e), \rho_r(e)$  are the total flow function and capacity constraint from  $G_{PW} - V(H_{pw})$  to  $H_{pw}$ . Moreover,  $W \leq \rho_j(e)$ .

**Proof:** I. Since  $\omega(e) = \sum_{i=0} \omega_{i+1}, \rho(e) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j \in G_{PW}$ , then the total flow

function and capacity constraint from  $H_{pw}$  or  $G_{PW} - V(H_{pw})$  to  $G_{PW}$  equal the total flow

function and capacity constraint from  $G_{PW}$  to  $H_{pw}$  or  $G_{PW} - V(H_{pw})$ , and

$H_{pw} + G_{PW} - V(H_{pw}) = G_{PW}$ , then  $PW = \omega_j(e) - \omega_r(e), PW = \rho_j(e) - \rho_r(e)$

Moreover,  $\omega_j(e) - \omega_r(e) \leq \omega_j(e) \leq \rho_j(e)$ .

**Definition (1.7):** Let  $\hat{G}_{PW} = (V, E^*(\hat{G}_{PW}), \Phi), \check{G}_{PW} = (V, E^*(\check{G}_{PW}), \Phi(\check{G}_{PW})) \subseteq G_{PW}$

be sub - Graph -  $\hat{PW}$  and sub - Graph -  $\check{PW}$  of network Graph - PW with

$V = \{v_i: i = 1, 2, 3, \dots, n\}$ , and

$E^*(\hat{R}_{PW}) = \{e_{ij} = v_i v_j: \deg v_i = \deg v_j \forall i, j \wedge n(e_{ij}) = v_i * v_j = \hat{PW}\}$ , and

$E^*(\check{R}_{PW}) = \{e_{ij} = v_i v_j: \deg v_i = \deg v_j \forall i, j \wedge n(e_{ij}) = v_i * v_j = \check{PW} \notin E^*(\hat{G}_{PW})\}$ .

Then  $\hat{G}_{PW}$  and  $\check{G}_{PW}$  are called **Regulars sub-graph- $\hat{PW}$**  and **sub - Graph -  $\check{PW}$**  denoted by

$\hat{R}_{PW}$  and  $\check{R}_{PW}$ .  $\hat{G}_{PW}$  and  $\check{G}_{PW}$  are **semi - reglars** if one or two vertices have the equal degree different of all vertices, that is,  $(\deg v_1 = \deg v_2) \neq (\deg v_i = \deg v_j) \forall i, j \wedge n(e_{ij}) = v_i * v_j = \hat{PW} \vee \check{PW}$ .

**Theorem (1.5):**  $\hat{R}_{PW} = (V, E^*, \Phi)$  is regular sub - graph -  $\hat{PW}$  iff  $\hat{R}_{PW}$  or  $\hat{R}_{PW} \cup \check{R}_{PW}$  are regulars.

Moreover, if  $\hat{R}_{PW}$  is semi - reglar sub - Graph -  $\hat{PW}$ , then  $\hat{R}_{PW} \vee \check{R}_{PW} \cup \check{R}_{PW}$  are semi - reglars or regulars.

Therefore,  $\hat{R}_{PW} = \check{R}_{PW}$  iff  $\hat{PW} = \check{PW}$ .

**Proof:** See definition (1,7).

**Definition (1.8):** Let  $\hat{G}_{PW}, \check{G}_{PW} \subseteq G_{PW}$  be Graph -  $\hat{PW}$  and Graph -  $\check{PW}$  of network

Graph - PW with  $V(\hat{G}_{PW}) = V(\check{G}_{PW}) = V(G_{PW}), E(\hat{G}_{PW}) \cap E(\check{G}_{PW}) = \emptyset$ , and

$\hat{G}_{PW} \cup \check{G}_{PW} = G_{PW}$ . Then  $\bar{\hat{G}}_{PW} = (V(G_{PW}), E^*(\bar{\hat{G}}_{PW}), \bar{\Phi})$  where

$E^*(\bar{\hat{G}}_{PW}) = E^*(\hat{G}_{PW}) \cup \{\bar{n}(e_{ij}): \bar{n}(e_{ij}) = \hat{PW} - n(e_{ij}), n(e_{ij}) \in E^*(\hat{G}_{PW})\}$ , and

$\bar{\check{G}}_{PW} = (V(G_{PW}), E^*(\bar{\check{G}}_{PW}), \bar{\Phi})$  where

$E^*(\bar{\check{G}}_{PW}) = E^*(\check{G}_{PW}) \cup \{\bar{n}(e_{ij}): \bar{n}(e_{ij}) = \check{PW} - n(e_{ij}), n(e_{ij}) \in E^*(\check{G}_{PW})\}$ .

**Theorem (1.6):**  $\bar{\hat{G}}_{PW} \cup \bar{\check{G}}_{PW} = (V(G_{PW}), E^*(\bar{\hat{G}}_{PW}) \cup E^*(\bar{\check{G}}_{PW}), \bar{\Phi} \cup \bar{\Phi})$ .

**Proof:** By definition (1.8) the result follows.

## A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW

**Theorem (1.7):** 1.  $\overline{\hat{G}_{PW} \cup \check{G}_{PW}} = \bar{G}_{PW} = \bar{\hat{G}}_{PW} \cap \bar{\check{G}}_{PW}$ , with  $\hat{G}_{PW} \cup \check{G}_{PW} = G_{PW}$   
 2.  $\bar{\hat{G}}_{PW} \cup \bar{\check{G}}_{PW} = G_{PW} = \overline{\hat{G}_{PW} \cap \check{G}_{PW}}$ , with  
 $E^*(\bar{\hat{G}}_{PW} \cup \bar{\check{G}}_{PW}) = \{\max(n(e_{ij}), \bar{n}(e_{ij})) : n(e_{ij}) \in E^*(\hat{G}_{PW}) \text{ or } E^*(\check{G}_{PW}) \text{ and } \bar{n}(e_{ij}) \in E^*(\bar{\hat{G}}_{PW}) \text{ or } E^*(\bar{\check{G}}_{PW})\}$ .

**Proof:** 1. Since  $\hat{G}_{PW} \cup \check{G}_{PW} = G_{PW}$ , then  $\overline{\hat{G}_{PW} \cup \check{G}_{PW}} = \bar{G}_{PW}$ , and by Morgan laws  $\overline{\hat{G}_{PW} \cup \check{G}_{PW}} = \bar{\hat{G}}_{PW} \cap \bar{\check{G}}_{PW}$ , and

2.  $\bar{\hat{G}}_{PW} \cup \bar{\check{G}}_{PW} = G_{PW} = \overline{\hat{G}_{PW} \cap \check{G}_{PW}}$  (by Morgan laws).

**Definition (1.9):** Let  $\dot{H}_{PW}, \check{H}_{PW} \subseteq G_{PW}$  be Graph –  $\dot{P}\dot{W}$  and Graph –  $\check{P}\check{W}$  network

Graph –  $PW$  with  $V(\dot{H}_{PW}) = V(\check{H}_{PW}), E(\dot{H}_{PW}) \cap E(\check{H}_{PW}) = \varphi$ , and

$\dot{H}_{PW} \cup \check{H}_{PW} = G_{PW}$ . Then  $\dot{H}_{PW} = (V(G_{PW}), E^*(\dot{H}_{PW}), \Phi(\dot{H}_{PW}))$  where

$$E^*(\dot{H}_{PW}) = E^*(\check{H}_{PW}) \cup \{\bar{n}(e_{ij}) : \bar{n}(e_{ij}) = \max(\dot{P}\dot{W}, \check{P}\check{W}) - n(e_{ij}), n(e_{ij}) \in E^*(\dot{H}_{PW})\}$$

and  $\check{H}_{PW} = (V(G_{PW}), E^*(\check{H}_{PW}), \Phi(\check{H}_{PW}))$  where

$$E^*(\check{H}_{PW}) = E^*(\dot{H}_{PW}) \cup \{\bar{n}(e_{ij}) : \bar{n}(e_{ij}) = \max(\dot{P}\dot{W}, \check{P}\check{W}) - n(e_{ij}), n(e_{ij}) \in E^*(\check{H}_{PW})\}$$

**Theorem (1.8):**  $\dot{H}_{PW} \cup \check{H}_{PW} = G_{PW} \cup \bar{G}_{PW}$ .

**Proof:** By definition (1.3) and (1.9), the result follows.

**Definition (1.10):** Let  $\dot{Y}_{PW}, \check{Y}_{PW} \subseteq G_{PW}$  be Graph –  $\dot{P}\dot{W}$  and Graph –  $\check{P}\check{W}$  of network

Graph –  $PW$  with  $\dot{Y}_{PW} = (V(G_{PW}), E^*(\dot{Y}_{PW}), \Phi(\dot{Y}_{PW}))$ ,

$E^*(\dot{Y}_{PW}) = \{\bar{n}(e_{ij}) : \bar{n}(e_{ij}) = \dot{P}\dot{W} - n(e_{ij}) \forall v_i, v_j \in V(G_{PW})\}$  and

$$\dot{Y}_{PW} = (V(G_{PW}), E^*(\dot{Y}_{PW}), \Phi(\dot{Y}_{PW})),$$

$E^*(\check{Y}_{PW}) = \{\bar{n}(e_{ij}) : \bar{n}(e_{ij}) = \check{P}\check{W} - n(e_{ij}) \forall v_i, v_j \in V(G_{PW})\}$ .

**Theorem (1.9):**  $\dot{G}_{PW} \cup \dot{Y}_{PW}$  and  $\check{G}_{PW} \cup \check{Y}_{PW}$  are complete Graph –  $\dot{P}\dot{W}$  and Graph –  $\check{P}\check{W}$ .

**Proof:** By definition (1.3) and (1.10)  $\dot{G}_{PW} \cup \bar{\dot{G}}_{PW} = \dot{G}_{PW} \cup \dot{Y}_{PW}$  or  $\check{G}_{PW} \cup \bar{\check{G}}_{PW} = \check{G}_{PW} \cup \check{Y}_{PW}$  with  $PW = \dot{P}\dot{W}$  and  $PW < \check{P}\check{W}$  or  $PW = \check{P}\check{W}$  and  $PW < \dot{P}\dot{W}$  which are completes.

**Definition (1.11):** Let  $\hat{G}_{PW} = (V, E^*, \Phi) \subseteq G_{PW}$  be connected Graph –  $\dot{P}\dot{W}$  with loops and

$V(\hat{G}_{PW}) = V(G_{PW})$  and let  $\hat{G}_{PW}^2 = (V, E_2^*, \Phi_2)$  define as-well-as  $\hat{G}_{PW}$  with joint two vertices  $v_i, v_j$  non-adjacency by  $n(e_{ij}) = v_i * v_j$  with  $n(e_{ir})$  and  $n(e_{rj})$  adjacency by  $v_r$ , that is, equivalence  $1 \leq d(v_i, v_j) \leq 2$ , where  $v_i$  is adjacent to  $v_j$  in  $\hat{G}_{PW}$ ,  $d(v_i, v_j)$  is the minimum numbers of vertices except first or last vertex between  $v_i, v_j$ . We can continuous to define  $\hat{G}_{PW}^3 = (V, E_3^*, \Phi_3), \dots$ , then,  $\hat{G}_{PW}^r = (V, E_r^*, \Phi_r)$  is called a **Power Graph- $\dot{P}\dot{W}$** .

**Theorem (1.10):**  $\hat{G}_{PW}^r = G_{PW}$  iff  $r = \max d(v_i, v_j)$

**Proof:** Since  $\hat{G}_{PW} = (V, E^*, \Phi) \subseteq G_{PW}$  with the loops, then  $\hat{G}_{PW}^2 \subseteq G_{PW}$ , and so on  $\hat{G}_{PW}^{r-1} \subseteq G_{PW}$ , and  $\hat{G}_{PW}^r = G_{PW}$  iff  $r = \max d(v_i, v_j)$ .

Now,  $\dot{R}_{PW}, \bar{\dot{G}}_{PW}, \bar{\check{G}}_{PW}, \dot{H}_{PW}, \check{H}_{PW}, \dot{Y}_{PW}, \check{Y}_{PW}$ , and  $\hat{G}_{PW}^r$  well define connected network Graph- $\dot{P}\dot{W}$  and Graph- $\check{P}\check{W}$ .

**Definition (1.12)** A **network symmetric Directed Graph- $PW$**  is triple  $SD_{PW} = (V, A^*, \Phi)$  consists of a non-empty Set  $V = \{v_1, v_2, \dots, v_n\}$  of objects called vertices, points, nodes, or just dots, together with directed pairs set of vertices

$$A^* = \left\{ \begin{array}{l} a = (a_{ij} = (v_i, v_j) \cup (v_j, v_i) = a_{ji}) : i \neq j \text{ or } i = j \wedge \\ n(a_{ij}) = v_i * v_j = v_j * v_i = \rho\omega \end{array} \right\}$$

is called edges, arcs or lines, satisfy  $PW = \max(n(a) = v_i * v_j = \rho\omega)$ , where  $P$  (non-negative integer) is the maximum numbers of  $\rho$ -arcs or loops between any pairs of vertices and  $W$  is the uniform weighted with  $PW = \max \rho\omega$  depended on  $\omega$  uniform sign weighted in  $\omega$  in  $A^*$ , and incident function  $\Phi$  from  $A$  to the set  $V^2$ , that is,

$\Phi: A^* \mapsto V^2$ . The **adjacency function matrix  $Am(SD_{PW})$**  define as

## A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW

$$Am(SD_{PW}) = \begin{cases} n(a) & \text{if } P > 0 \wedge w > 0 \wedge v_i \text{ joined to } v_j; \\ 0 \text{ or } \infty & \text{if } \rho\omega = 0 ; \\ -n(a) & \text{if } P > 0 \wedge w < 0 \wedge v_i \text{ joined to } v_j. \end{cases}$$

clear that it is if  $P = W = 1$ , then the symmetric Digraph-  $PW$  is natural **symmetric Digraph-1**, and if  $P > 0, W = 1$  the symmetric Digraph-  $PW$  is **mullet symmetric Digraph- P**, and if  $P = 1, W \in \mathbb{R}^+$  the symmetric Digraph-  $PW$  is **weighted symmetric Digraph- W** and  $\Phi$  in symmetric Digraph-  $PW$  satisfy

$\forall a \in A^* \exists \{(v_i, v_j), (v_j, v_i)\} \in P_2(V) \ni \Phi(a) = \{(v_i, v_j), (v_j, v_i)\}$ ,  $\Phi$  is one -to-one if and only if  $\Phi(a_i) = \Phi(a_j) \Rightarrow a_i = a_j$  or  $a_i \neq a_j \Rightarrow \Phi(a_i) \neq \Phi(a_j)$ , onto if  $\forall \{(v_i, v_j), (v_j, v_i)\} \in P_2(V) \exists a \in A^* \ni \Phi(\{(v_i, v_j), (v_j, v_i)\}) = a$  or  $\Phi(P_2(V)) = A^*$ , and if  $\Phi$  one-to-one and onto is called corresponding then there is an inverse  $\Phi^{-1}$  of  $\Phi$ .

In the following definitions and theorems, incident function and adjacency function matrix define as well as.

**Definition (1.13):** In the **network symmetric Digraph- PW**, the weight  $\omega$  is called **symmetric flow function** for each arcs of  $\rho \in n(a)$  with  $\omega_i \leq \omega_{i+1}$  and  $\rho = \sum_{i=1}^{\rho} \omega_i$  between any adjacent two pair vertices is called **symmetric capacity constraint** with  $\omega \leq \rho = \sum_{i=1}^{\rho} \omega_i$  a and the total symmetric flow function for any vertex is

$\omega(a) = \sum_{i=0} \omega_{i+1}$ , and the total symmetric capacity constraint for any vertex is  $\rho(a) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j$ , with  $\omega(a) \leq \rho(a)$ .

If  $\omega(a) = \sum_{i=0} \omega_{i+1} = PW$  then  $PW$  is called **value symmetric flow**, therefore, for any intermediate vertex if the total symmetric flow is  $PW$  of these vertices is called **symmetric flow conservation**.

By these arguments the name **symmetric Digraph- PW** and **network symmetric Digraph- PW** are the same.

**Definition (1.14):** The **complement** of  $SD_{PW}$  is  $S\bar{D}_{PW} = (V, \bar{A}^*, \bar{\Phi})$  with

$V(SD_{PW}) = V(S\bar{D}_{PW})$ , and

$\bar{A}^* = \{a = (v_i v_j) \cup (v_j v_i) \wedge \bar{n}(a) = PW - n(a) \forall v_i, v_j \in V(SD_{PW})\}$ , and

$\bar{\Phi}: \bar{A}^* \rightarrow V^2$ . Therefore  $SD_{PW} \cup S\bar{D}_{PW} = K_{|V|} - PW$  is called **complete symmetric Digraph- PW**.

**Definition (1.15):** The **regular** of  $SD_{PW}$  is  $SR_{PW} = (V, A^*, \Phi)$  with

$A^* = \{a = (v_i v_j) \cup (v_j v_i): \deg v_i = \deg v_j \wedge n(a) = v_i * v_j = \rho\omega = PW \forall v_i, v_j \in V(SD_{PW})\}$

and  $SD_{PW}$  is **sime – regular** if

$(indeg v_1 = indeg v_2) \neq (indeg v_i = indeg v_j) \forall i, j \wedge n(a) = v_i * v_j = PW$

**Theorem (1.12):** If a symmetric Digraph-  $PW$   $SD_{PW} = (V, A^*, \Phi)$  or  $S\bar{D}_{PW}$  has an isolated vertex, then  $SD_{PW} - v$  or  $SD_{PW} \cup S\bar{D}_{PW}$  has not a isolated vertex.

**Proof:** Since  $SD_{PW} = (V, A^*, \Phi)$  is defined on operations, then the operation module has only isolated vertex, and the operation of product has isolated vertex, if one of vertex is zero, and the anther operations  $S\bar{D}_{PW}$  has isolated vertex, therefore  $SD_{PW} - v$  or

$SD_{PW} \cup S\bar{D}_{PW}$  has not an isolated vertex.

Now, by these theorem  $SD_{PW}, S\bar{D}_{PW}$  and  $SD_{PW} \cup \bar{G}_{PW}$  well define are connected and we can be taken these network.

**Definition (1.16):** In the value flow of the networks, if  $\sum_{i=0} \omega_{i+1} > PW$ , then there is  $\sum_{j=0} \omega_{j+1}$  such that  $\sum_{i=0} \omega_{i+1} - \sum_{j=0} \omega_{j+1} = PW$ , and  $PW \leq \rho(a) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j$ .

**Definition (1.17):** A **sub-symmetric Digraphs-  $\dot{P}W$**  of a symmetric Digraph-  $PW$  is  $SH_{pw}$  or  $S\dot{D}_{PW} = (V, A^*, \Phi) \subseteq SD_{PW}$  with  $V(SH_{pw}) \subseteq V(SD_{PW})$  or

$V(S\dot{D}_{PW}) = V(SD_{PW}), A^*(SH_{pw})$  or

$A^*(S\dot{D}_{PW}) \subseteq A^*(SD_{PW})$ , and  $\Phi(SH_{pw})$  or  $\Phi(S\dot{D}_{PW}) \subseteq \Phi(SD_{PW})$ .

Now, let  $SH_{pw} = (V(SH), A^*(SH), \Phi(SH))$  and

$S\dot{Y}_{pw} = (V(SY), A^*(SY), \Phi(SY)) \subseteq SD_{PW} = (V, A^*, \Phi)$  with

$$V(SY) = V(SD_{PW}) - V(SH),$$

## A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW

$A^*(Y) = \{a = (v_i, v_j) \cup (v_j, v_i) : n(a) = v_i * v_j = \rho\omega \notin A^*(SH)\}$ . Then

$SH_{pw} \cup SY_{pw} \subseteq SD_{PW}$ , and  $SH_{pw} + SY_{pw} = SD_{PW}$ , where  $V(SH) \cup V(SY) = V(SD_{PW})$  and  $A^*(SH) \cup A^*(SY) \cup \{a = (v_i, v_j) \cup (v_j, v_i) : v_i \in V(H), v_j \in V(Y) \forall v_i, v_j, n(a) = v_i * v_j = \rho\omega\}$ .

Let  $SH_{pw} = (V(SH), A^*(SH), \Phi(SH))$ ,

$SY_{pw} = (V(SY), A^*(SY), \Phi(SY)) \subseteq SD_{PW} = (V, A^*, \Phi) \subseteq SD_{PW}$  with

$$V(SY) = V(SD) - V(SH),$$

$A^*(Y) = \{a = (v_i, v_j) \cup (v_j, v_i) : n(a) = v_i * v_j = \rho\omega \notin A^*(SH)\}$ . Then

$SH_{pw} \cup SY_{pw} \subseteq SD_{PW}$ , and  $SH_{pw} \cup SY_{pw} = SD_{PW}$  iff

$$A^*(SH) \cup A^*(SY) \cup \{a = (v_i, v_j) \cup (v_j, v_i) : v_i \in V(SH), v_j \in V(SY) \wedge$$

$n(a) = v_i * v_j = \rho\omega \in A^*(SD_{PW})\}$ .

Now, we can give the following results on the network symmetric Digraph- PW on symmetric flow function, value symmetric flow and symmetric capacity constraint.

**Theorem (1.13):** If  $SH_{pw} = (V(SH_{pw}), A^*(SH_{pw}), \Phi(SH_{pw})) \subseteq SD_{PW} = (V, A^*, \Phi)$  of network symmetric Digraph- PW, with  $pw = PW$ , then

$$SH_{pw} + SD_{PW} - V(SH_{pw}) = SD_{PW}.$$

**Proof:** By definition (1.17) the result follows.

**Theorem (1.14):** Let  $\omega(a) = \sum_{i=0} \omega_{i+1}$ ,  $\rho(a) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j > PW$  be total symmetric flow function and symmetric capacity constraint of networks symmetric Digraph- PW. Then there is

$\sum_{r=0} \omega_{i+1}$ ,  $\rho(a) = \sum_{r=1} (\sum_{i=1}^{\rho} \omega_i)_r \in SH_{pw}$  or  $SD_{PW} - V(SH_{pw}) \subseteq SD_{PW} = (V, A^*, \Phi)$  such that

$$I. \omega(a) - \sum_{r=0} \omega_{r+1} = PW, \text{ and}$$

$$II. \rho(a) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j - \rho(a) = \sum_{r=1} (\sum_{i=1}^{\rho} \omega_i)_r = PW.$$

**Proof:** I. Since  $\sum_{r=0} \omega_{i+1} \in SH_{pw}$  or  $SD_{PW} - V(SH_{pw})$ , then  $\sum_{r=0} \omega_{r+1} \in SD_{PW}$ , and by definition (1.5) the result follows, and

II. Since  $\rho(a) = \sum_{r=1} (\sum_{i=1}^{\rho} \omega_i)_r \in SH_{pw}$  or  $SD_{PW} - V(SH_{pw})$ , then

$\rho(a) = \sum_{r=1} (\sum_{i=1}^{\rho} \omega_i)_r \in SD_{PW}$ , and the result follows.

**Theorem (1.15):** Let  $\omega(a) = \sum_{i=0} \omega_{i+1}$ ,  $\rho(a) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j$  be total symmetric flow function and symmetric capacity constraint of networks symmetric Digraph- PW have value symmetric flow and let,  $\omega_j(a), \rho_j(a)$  be total symmetric flow function and symmetric capacity constraint of sub- symmetric Digraph- pw from  $SH_{pw}$  to  $SD_{PW} - V(SH_{pw})$ . Then  $PW = \omega_j(a) - \omega_r(a)$ ,  $PW = \rho_j(a) - \rho_r(a)$ , where  $\omega_r(a)$   $\rho_r(a)$  are the total symmetric flow function and symmetric capacity constraint from  $G_{PW} - V(H_{pw})$  to  $H_{pw}$ . Moreover,  $\leq \rho_j(a)$ .

**Proof:** Since  $\omega(a) = \sum_{i=0} \omega_{i+1}$ ,  $\rho(a) = \sum_{j=1} (\sum_{i=1}^{\rho} \omega_i)_j \in G_{PW}$ , then the total

symmetric flow function and symmetric capacity constraint from  $SH_{pw}$  or

$SD_{PW} - V(SH_{pw})$  to  $SD_{PW}$  equal the total symmetric flow function and

symmetric capacity constraint from  $SD_{PW}$  to  $SH_{pw}$  or  $SD_{PW} - V(SH_{pw})$ , and

$SH_{pw} + SD_{PW} - V(SH_{pw}) = SD_{PW}$ , then  $PW = \omega_j(a) - \omega_r(a)$ ,  $PW = \rho_j(a) - \rho_r(a)$

Moreover,  $\omega_j(a) - \omega_r(a) \leq \omega_j(a) \leq \rho_j(a)$ .

**Definition (1.18):** Let  $SD_{PW} = (V, A^*(SD_{PW}), \Phi)$  and

$$SD_{PW} = (V, A^*(SD_{PW}), \Phi(SD_{PW})) \subseteq SD_{PW}$$

be sub - symmetric Digraph - PW and sub - symmetric Digraph - PW of network

symmetric Digraph - PW with  $V = \{v_i : i = 1, 2, 3, \dots, n\}$ , and



**A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW**

$$A^*(S\check{R}_{PW}) = \{a = v_i v_j : deg v_i = deg v_j \forall i, j \wedge n(a) = v_i * v_j = \check{P}\check{W}\}, \text{ and}$$

$$A^*(S\check{R}_{PW}) = \{a = v_i v_j : indeg v_i = outdeg v_j \forall i, j \wedge$$

$$n(a) = v_i * v_j = \check{P}\check{W} \notin E^*(S\check{D}_{PW})\}.$$

Then  $S\check{D}_{PW}$  and  $S\check{D}_{PW}$  are called **Regulars sub- symmetric Digraph- $\check{P}\check{W}$**  and **sub – symmetric Digraph –  $\check{P}\check{W}$**  denoted by  $S\check{R}_{PW}$  and  $S\check{R}_{PW}$ .  $S\check{D}_{PW}$  and  $S\check{D}_{PW}$  are **semi – regular** if one or two vertices have the equal degree different of all vertices, that is ,

$$(indeg v_1 = indeg v_2) \neq (indeg v_i = indeg v_j) \forall i, j \wedge n(a) = v_i * v_j = \check{P}\check{W} \vee \check{P}\check{W}.$$

**Theorem (1.15):**  $\check{R}_{PW} = (V, A^*, \Phi)$  is regular sub – symmetric Digraph –  $\check{P}\check{W}$  iff  $S\check{R}_{PW}$

or  $S\check{R}_{PW} \cup S\check{R}_{PW}$  are regulars. Moreover, if  $S\check{R}_{PW}$  is semi – regular sub – symmetric Digraph –  $\check{P}\check{W}$ , then  $S\check{R}_{PW}$  or  $S\check{R}_{PW} \cup S\check{R}_{PW}$  are semi – regulars or regulars. Therefore,  $S\check{R}_{PW} = S\check{R}_{PW}$  iff  $\check{P}\check{W} = \check{P}\check{W}$ .

**Proof:** See definition (1,7).

**Definition (1.19):** Let  $S\check{D}_{PW}, S\check{D}_{PW} \subseteq SD_{PW}$  be Graph –  $\check{P}\check{W}$  and Graph –  $\check{P}\check{W}$  of network

$$\text{Graph – PW with } V(S\check{D}_{PW}) = V(S\check{D}_{PW}) = V(SD_{PW}), A^*(S\check{D}_{PW}) \cap A^*(S\check{D}_{PW}) = \varphi, \text{ and}$$

$$S\check{D}_{PW} \cup S\check{D}_{PW} = SD_{PW}. \text{ Then } S\check{D}_{PW} = (V(SD_{PW}), A^*(S\check{D}_{PW}), \check{\Phi}) \text{ where}$$

$$A^*(S\check{D}_{PW}) = A^*(S\check{D}_{PW}) \cup \{\bar{n}(a) : \bar{n}(a) = \check{P}\check{W} - n(a), n(a) \in A^*(S\check{D}_{PW})\}, \text{ and}$$

$$S\check{D}_{PW} = (V(SD_{PW}), A^*(S\check{D}_{PW}), \check{\Phi}) \text{ where}$$

$$A^*(S\check{D}_{PW}) = A^*(S\check{D}_{PW}) \cup \{\bar{n}(a) : \bar{n}(a) = \check{P}\check{W} - n(a), n(a) \in A^*(S\check{D}_{PW})\}.$$

**Theorem (1.16):**  $S\check{D}_{PW} \cup S\check{D}_{PW} = (V(SD_{PW}), A^*(S\check{D}_{PW}) \cup A^*(S\check{D}_{PW}), \check{\Phi} \cup \check{\Phi})$ .

**Proof:** By definition (1.8) the result follows.

**Theorem (1.17):** 1.  $\overline{S\check{D}_{PW} \cup S\check{D}_{PW}} = S\check{D}_{PW} = S\check{D}_{PW} \cap S\check{D}_{PW}$ , with  $S\check{D}_{PW} \cup S\check{D}_{PW} = SD_{PW}$

$$2. S\check{D}_{PW} \cup S\check{D}_{PW} = SD_{PW} = \overline{S\check{D}_{PW} \cap S\check{D}_{PW}}, \text{ with}$$

$$A^*(S\check{D}_{PW} \cup S\check{D}_{PW}) = \{\max(n(a), \bar{n}(a)) : n(a) \in A^*(S\check{D}_{PW}) \text{ or } A^*(S\check{D}_{PW}) \text{ and}$$

$$\bar{n}(a) \in A^*(S\check{D}_{PW}) \text{ or } A^*(S\check{D}_{PW})\}.$$

**Proof:** 1. Since  $S\check{D}_{PW} \cup S\check{D}_{PW} = SD_{PW}$ , then  $\overline{S\check{D}_{PW} \cup S\check{D}_{PW}} = S\check{D}_{PW} = S\check{D}_{PW} \cap S\check{D}_{PW}$

(by De-Morgan laws), and so 2.

**Definition (1.20):** Let  $S\check{H}_{PW}, S\check{H}_{PW} \subseteq SD_{PW}$  be sub – symmetric Digraph –  $\check{P}\check{W}$  and sub –

symmetric Digraph –  $\check{P}\check{W}$  of network symmetric Digraph – PW with

$$V(S\check{H}_{PW}) = V(S\check{H}_{PW}), E(S\check{H}_{PW}) \cap E(S\check{H}_{PW}) = \varphi, \text{ and}$$

$$S\check{H}_{PW} \cup S\check{H}_{PW} = SD_{PW}. \text{ Then } S\check{H}_{PW} = (V(SD_{PW}), A^*(S\check{H}_{PW}), \Phi(S\check{H}_{PW})) \text{ where}$$

$$A^*(S\check{H}_{PW}) = A^*(S\check{H}_{PW}) \cup \{\bar{n}(a) : \bar{n}(a) = \max(\check{P}\check{W}, \check{P}\check{W}) - n(a), n(a) \in A^*(S\check{H}_{PW})\} \text{ and}$$

$$S\check{H}_{PW} = (V(SD_{PW}), A^*(S\check{H}_{PW}), \Phi(S\check{H}_{PW})) \text{ where}$$

$$A^*(S\check{H}_{PW}) = A^*(S\check{H}_{PW}) \cup \{\bar{n}(a) : \bar{n}(a) = \max(\check{P}\check{W}, \check{P}\check{W}) - n(a), n(a) \in A^*(S\check{H}_{PW})\}$$

**Theorem (1.18):**  $S\check{H}_{PW} \cup S\check{H}_{PW} = SD_{PW} \cup S\check{D}_{PW}$ .

**Proof:** By definition (1.3) and (1.9), the result follows.

## A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW

**Definition (1.21):** Let  $S\check{Y}_{PW}, S\check{Y}_{PW} \subseteq SD_{PW}$  be sub – symmetric Digraph –  $\check{P}\check{W}$  and sub – symmetric Digraph –  $\check{P}\check{W}$  of network

symmetric Digraph – PW with  $S\check{Y}_{PW} = (V(SD_{PW}), A^*(S\check{Y}_{PW}), \Phi(S\check{Y}_{PW}))$ ,

$A^*(S\check{Y}_{PW}) = \{\bar{n}(a): \bar{n}(a) = \check{P}\check{W} - n(a) \forall v_i, v_j \in V(SD_{PW})\}$  and

$S\check{Y}_{PW} = (V(SD_{PW}), A^*(S\check{Y}_{PW}), \Phi(S\check{Y}_{PW}))$ ,

$A^*(S\check{Y}_{PW}) = \{\bar{n}(a): \bar{n}(a) = \check{P}\check{W} - n(a) \forall v_i, v_j \in V(SD_{PW})\}$ .

**Theorem (1.19):**  $S\check{D}_{PW} \cup S\check{Y}_{PW}$  and  $S\check{D}_{PW} \cup S\check{Y}_{PW}$  are complete symmetric Digraph –  $\check{P}\check{W}$  and symmetric Digraph- $\check{P}\check{W}$ .

**Proof:** By definition (1.3) and (1.10) we have

$SD_{PW} \cup S\check{D}_{PW} = S\check{D}_{PW} \cup S\check{Y}_{PW}$  or  $S\check{D}_{PW} \cup S\check{Y}_{PW}$  with

$PW = \check{P}\check{W}$  and  $PW < \check{P}\check{W}$  or  $PW = \check{P}\check{W}$  and  $PW < \check{P}\check{W}$  which are completes.

**Definition (1.22):** Let  $S\check{D}_{PW} = (V, A^*, \Phi) \subseteq SD_{PW}$  be connected sub – symmetric

Digraph –  $\check{P}\check{W}$  with loops and  $V(S\check{D}_{PW}) = V(SD_{PW})$  and let  $S\check{D}_{PW}^2 = (V, A_2^*, \Phi_2)$  define as-well-as  $S\check{D}_{PW}$  with joint two vertices  $v_i, v_j$  non-adjacency by  $n(a_{ij}) = v_i * v_j$  with  $n(a_{ir})$  and  $n(a_{rj})$  adjacency by  $v_r$ , that is, equivalence  $1 \leq d(v_i, v_j) \leq 2$ , where  $v_i$  is adjacent to  $v_j$  in  $S\check{D}_{PW}$ ,  $d(v_i, v_j)$  is the minimum numbers of vertices except first or last vertex between  $v_i, v_j$ . We

can continuous define  $S\check{D}_{PW}^3 = (V, A_3^*, \Phi_3), \dots$ , then,  $S\check{D}_{PW}^r = (V, A_r^*, \Phi_r)$  is called a **Power symmetric Digraph- $\check{P}\check{W}$** .

**Theorem (1.20):**  $S\check{D}_{PW}^r = SD_{PW}$  iff  $r = \max d(v_i, v_j)$

**Proof:** Since  $S\check{D}_{PW} = (V, A^*, \Phi) \subseteq SD_{PW}$  with the loops, then  $S\check{D}_{PW}^2 \subseteq SD_{PW}$ , and so on

$S\check{D}_{PW}^{r-1} \subseteq SD_{PW}$ , and  $S\check{D}_{PW}^r = SD_{PW}$  iff  $r = \max d(v_i, v_j)$ .

Now,  $S\check{R}_{PW}, S\check{D}_{PW}, S\check{D}_{PW}, S\check{H}_{PW}, S\check{H}_{PW}, S\check{Y}_{PW}, S\check{Y}_{PW}$ , and  $S\check{D}_{PW}^r$  well define connected network symmetric Digraph- $\check{P}\check{W}$  and symmetric Digraph- $\check{P}\check{W}$ .

**Theorem (1.21):**  $G_{PW} = SD_{PW}$  iff  $V(G_{PW}) = V(SD_{PW})$ , and  $D_{PW} = D_{PW}^{-1}$

**Proof:**  $G_{PW} = SD_{PW}$  iff  $SD_{PW} = D_{PW} \cup D_{PW}^{-1}$  satisfy  $V(G_{PW}) = V(SD_{PW})$ ,

$e = \{v_i, v_j\} = a = (a_{ij} = (v_i, v_j) \cup (v_j, v_i) = a_{ij}^{-1}) \forall i \neq j, i = j$ , and

$\deg_{v_i \in V(G_{PW})} v_i = \text{indeg}_{v_i \in V(SD_{PW})} v_i = \text{oudeg}_{v_i \in V(SD_{PW})} v_i$ , and

$\deg_{n(e_{ij}) \in E^*} n(e_{ij}) = \text{indeg}_{n(a) \in A^*} n(a) = \text{oudeg}_{n(a) \in A^*} n(a)$

Moreover  $E^* = A^*$  iff  $E^* = (A(D_{PW}) = A(D_{PW}^{-1}))$ , and  $P_2(V) = V^2$  iff  $G_{PW} = (D_{PW} = D_{PW}^{-1})$ .

### 2. Networks Change Graph-PW and Change Symmetric Digraph-PW:

**Definition (2.1):** Let  $G_{PW} = (V, E, \Phi)$  be Graph-PW. Then we can define the **change Graph-PW** as  $L(G_{PW}) = (V(E^*), E^*(L), L(\Phi))$  with  $V(E^*) = \{n(e_{ij}): n(e_{ij}) \in E^*(G_{PW})\}$ ,

$E^*(L) = \{e_{ij} = n(e_{ir})n(e_{rj}): n(e_{ij}) = n(e_{ir}) * n(e_{rj}) = \rho\omega\}$ ,

$E^*(L) = \{e_{ij} = n(e_{ir})n(e_{rj}): n(e_{ij}) = v_i * v_j = \rho\omega\}$ ,

$E^*(L) = \{e_{ij} = n(e_{ir})n(e_{rj}): n(e_{ij}) = v_r * v_r = \rho\omega\}$  or

$E^*(L) = \{e_{ij} = n(e_{ir})n(e_{rj}): n(e_{ij}) = v_r = \rho\omega\}$ , that is,

the set  $n(e_{ij})$ -edges of  $E^*(G_{PW})$  is vertices in  $L(G_{PW})$ , and the  $e_{ij}$  in  $E^*(L)$  if and only if  $n(e_{ir})$  and  $n(e_{rj})$  are adjacency of vertex  $v_r$ ,

$L(\Phi): E^*(L) \rightarrow P_2(V(E^*))$ , and

$PW = \max \rho\omega = \max n(e_{ij}) * n(e_{rj}), v_i * v_j, v_r * v_r$  or  $v_r = \rho\omega$ .



## A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW

Moreover, we can be found  $L^2(G_{PW}) = L(L(G_{PW})), \dots, L^n(G_{PW}) = L(L^{n-1}(G_{PW}))$ .  
Applied all definitions in the definition(1.1) and definition (1.2) in the definition (2.1).

**Definition (2.2):** The complement of network change Graph-PW

$L(G_{PW}) = (V(E^*), E^*(L), L(\Phi))$  is  $\bar{L}(G_{PW}) = (V(E^*), \bar{E}^*(L), \bar{L}(\Phi))$  with

$\bar{E}^*(L) = \{e_{ij} : \bar{n}(e_{ij}) = PW - n(e_{ij}), i \neq j \text{ or } i = j\}$

$\bar{L}(\Phi) : \bar{E}^*(L) \rightarrow P_2(V(E^*))$ , and

$$PW = \max \bar{n}(e_{ij}).$$

Moreover,  $L(G_{PW}) \cup \bar{E}^*(L) = K_{|V(E^*)|} - PW$ .

**Definition (2.3):** Let  $G_{PW} = (V, E^*, \Phi)$  be network regular Graph-PW. Then we can define the network change regular Graph-PW as  $L(G_{PW}) = (V(E^*), E^*(L), \Phi(L))$  with

$V(E^*) = \{PW : PW \in E^*(G_{PW})\}$ ,

$$E^*(L) = \left\{ e_{ij} : n(e_{ij}) = \frac{|V|^2(|V| - 1)}{2} PW * PW = \rho\omega \right\},$$

$\Phi(L) : E^*(L) \rightarrow P_2(V(E^*))$ , and  $PW = \max \rho\omega$ .

**Definition (2.4):** Applied definitions (1.5) and (1.6) in the network change Graph-PW.

**Theorem (2.1):** Applied theorems (1.1),(1.2),(1.3) and (1.4) in the network change Graph-PW.

Clear that if  $G_{PW} = (V, E^*, \Phi)$  is connected Graph-PW. Then  $L(G_{PW}) = (V(E^*), E^*(L), L(\Phi))$  is connected change Graph-PW, if  $G_{PW} = (V, E^*, \Phi)$  has a soiled vertex, then  $\bar{G}_{PW} = (V, \bar{E}^*, \bar{\Phi})$  or  $G_{PW} \cup \bar{G}_{PW}$  are connected, so  $L(\bar{G}_{PW})$  and  $L(G_{PW} \cup \bar{G}_{PW})$ , moreover can be looking the change Graph-PW  $L(G_{PW} \cup \bar{G}_{PW})$  has only one vertex with the loops. If  $G_{PW} = (V, E^*, \Phi)$  has a soiled vertex, then

$L(G_{PW}) = (V(E^*), E^*(L), L(\Phi))$  is connected.

**Definition (2.5):** The definitions (1.7) define similar of definition (2.3) in the network change Graph-PW.

**Definition (2.6):** Applied the definitions (1.8), (1.9),(1.10) and (1.11) in the network change Graph-PW.

**Theorem (2.2):** Applied theorems (1.6), (1.7),(1.8), (1.9) and (1.10) in the network change Graph-PW.

Now, let  $\hat{G}_{PW}, \check{G}_{PW} \subseteq G_{PW}$  be Graph -  $\hat{P}W$  and Graph -  $\check{P}W$  with

$V(\hat{G}_{PW}) = V(\check{G}_{PW}), E(\hat{G}_{PW}) \cap E(\check{G}_{PW}) = \varphi$ , and  $\hat{G}_{PW} \cup \check{G}_{PW} = G_{PW}$ .

Clear that  $\hat{G}_{PW}$  is  $L(\hat{G}_{PW})$ , and  $L^2(\hat{G}_{PW}) = L(L(\hat{G}_{PW})), \dots, L^n(\hat{G}_{PW}) = L(L^{n-1}(\hat{G}_{PW}))$  and  $\check{G}_{PW}$  is  $L(\check{G}_{PW})$ , and  $L^2(\check{G}_{PW}) = L(L(\check{G}_{PW})), \dots, L^n(\check{G}_{PW}) = L(L^{n-1}(\check{G}_{PW}))$ .

**Theorem (2.3):**  $\bar{\hat{G}}_{PW}, \bar{\check{G}}_{PW}, \hat{H}_{PW}$  and  $\check{H}_{PW}$  have network change Graph-PW and

$$L(\bar{\hat{G}}_{PW}), L^2(\bar{\hat{G}}_{PW}) = L(L(\bar{\hat{G}}_{PW})), \dots, L^n(\bar{\hat{G}}_{PW}) = L(L^{n-1}(\bar{\hat{G}}_{PW})),$$

$$L(\bar{\check{G}}_{PW}), L^2(\bar{\check{G}}_{PW}) = L(L(\bar{\check{G}}_{PW})), \dots, L^n(\bar{\check{G}}_{PW}) = L(L^{n-1}(\bar{\check{G}}_{PW})),$$

$L(\hat{H}_{PW}), L^2(\hat{H}_{PW}) = L(L(\hat{H}_{PW})), \dots, L^n(\hat{H}_{PW}) = L(L^{n-1}(\hat{H}_{PW}))$  and

$L(\check{H}_{PW}), L^2(\check{H}_{PW}) = L(L(\check{H}_{PW})), \dots, L^n(\check{H}_{PW}) = L(L^{n-1}(\check{H}_{PW}))$ .

**Proof:** By the definition (2.1) the result follows.

**Theorem (2.4):** Every  $\hat{G}_{PW}^r = (V, E_r^*, \Phi_r)$  there is  $L(\hat{G}_{PW}^r), r = 1, 2, 3, \dots, \max d(v_i, v_j)$ .

Moreover,  $L^2(\hat{G}_{PW}^r) = L(L(\hat{G}_{PW}^r)), \dots, L^n(\hat{G}_{PW}^r) = L^{n-1}(\hat{G}_{PW}^r)$ .

# A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW

**Proof:** By theorem (1.10) and the definition (2.1) the result follows.

**Definition (2.7):** Let  $SD_{PW} = (V, A^*, \Phi)$  be Symmetric Digraph-PW. Then we can define the change symmetric Digraph-PW is

$L(SD_{PW}) = (V(A^*), A^*(L), L(\Phi))$  with

$V(A^*) = \{n(a) : a \in A^*(SD_{PW})\}$ ,

$A^*(L) = \{a_{ij} = n(a_{ir}) n(a_{rj}) : n(a_{ij}) = n(a_{ir}) * n(a_{rj}) = \rho\omega\}$ ,

$A^*(L) = \{a_{ij} = n(a_{ir}) n(a_{rj}) : n(a_{ij}) = v_i * v_j = \rho\omega\}$ ,

$A^*(L) = \{a_{ij} = n(a_{ir}) n(a_{rj}) : n(a_{ij}) = v_r * v_r = \rho\omega\}$  or

$A^*(L) = \{a_{ij} = n(a_{ir}) n(a_{rj}) : n(a_{ij}) = v_r = \rho\omega\}$ ,

that is, the symmetric set  $n(a_{ij})$ -arcs of  $A^*(SD_{PW})$  is vertices in  $L(SD_{PW})$ , and the  $a_{ij}$  in  $A^*(L)$  if and only if  $n(a_{ir})$  and  $n(a_{rj})$  are adjacency of vertex  $v_r$ ,

$$L(\Phi): A^*(L) \rightarrow V^2(A^*), \text{ and}$$

$$PW = \max \rho\omega = \max n(a_{ij}) * n(a_{rj}), v_i * v_j, v_r * v_r \text{ or } v_r = \rho\omega.$$

Moreover, we can be found  $L^2(SD_{PW}) = L(L(SD_{PW})), \dots, L^n(SD_{PW}) = L(L^{n-1}(SD_{PW}))$ .

Applied all definitions in the definition (1.12) and definition (1.13) in the definition (2.7).

**Definition (2.8):** The complement of change symmetric Digraph-PW

$L(SD_{PW}) = (V(A^*), A^*(L), L(\Phi))$  is  $\bar{L}(SD_{PW}) = (V(A^*), \bar{A}^*(\bar{L}), \bar{L}(\bar{\Phi}))$  with

$\bar{A}^*(\bar{L}) = \{a_{ij} : \bar{n}(a_{ij}) = PW - n(a_{ij}), i \neq j \text{ or } i = j\}$

$\bar{L}(\bar{\Phi}): \bar{A}^*(\bar{L}) \rightarrow V^2(A^*)$ , and

$$PW = \max \bar{n}(a_{ij}) = \max \rho\omega.$$

Moreover,  $L(SD_{PW}) \cup \bar{A}^*(\bar{L}) = K_{|V(A^*)|} - PW$

**Definition (2.9):** Let  $SD_{PW} = (V, A^*, \Phi)$  be regular symmetric Digraph-PW. Then we can define the change regular symmetric Digraph-PW is

$L(G_{PW}) = (V(E^*), E^*(L), \Phi(L))$  with  $V(E^*) = \{PW : PW \in E^*(G_{PW})\}$ ,

$$A^*(L) = \left\{ a_{ij} : n(a_{ij}) = \frac{|V|^2(|V| - 1)}{2} PW * PW = \rho\omega \right\}$$

$$\Phi(L): A^*(L) \rightarrow P_2(V(E^*)), \text{ and } PW = \max \rho\omega$$

**Definition (2.10):** Applied definitions (1.14) and (1.15) in the network change symmetric Digraph-PW.

**Theorem (2.5):** Applied theorems (1.12), (1.3), (1.14) and (1.15) in the network change symmetric Digraph-PW.

Clear that if  $SD_{PW} = (V, A^*, \Phi)$  is connected symmetric Digraph-PW, then

$L(SD_{PW}) = (V(A^*), A^*(L), L(\Phi))$  is connected change symmetric Digraph-PW, if  $SD_{PW} = (V, A^*, \Phi)$  has a soiled vertex, then  $S\bar{D}_{PW} = (V, \bar{E}^*, \bar{\Phi})$  or  $SD_{PW} \cup S\bar{D}_{PW}$  are connected, so  $L(S\bar{D}_{PW})$  and  $L(SD_{PW} \cup S\bar{D}_{PW})$ , moreover can be looking the change symmetric Digraph-PW  $L(SD_{PW} \cup S\bar{D}_{PW})$  has only one vertex with the loops. If  $SD_{PW} = (V, A^*, \Phi)$  has a soiled vertex, then  $L(SD_{PW}) = (V(A^*), A^*(L), L(\Phi))$  is connected.

**Definition (2.11):** The definition (1.16) define similar of definition (2.9) in the network change symmetric Digraph-PW.

**Definition (2.12):** Applied the definitions (1.17), (1.18), (1.19) and (1.20) in the network change symmetric Digraph-PW.

**Theorem (2.6):** Applied theorems (1.17), (1.18), (1.19), and (1.20) in the network change symmetric Digraph-PW.

Now, let  $\hat{G}_{PW}, \check{G}_{PW} \subseteq G_{PW}$  be Graph -  $\hat{P}W$  and Graph -  $\check{P}W$  with

$V(\hat{G}_{PW}) = V(\check{G}_{PW}), E(\hat{G}_{PW}) \cap E(\check{G}_{PW}) = \varphi$ , and  $\hat{G}_{PW} \cup \check{G}_{PW} = G_{PW}$ .

## A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW

Clear that  $\dot{G}_{PW}$  is  $L(\dot{G}_{PW})$ , and  $L^2(\dot{G}_{PW}) = L(L(\dot{G}_{PW}))$ , ...,  $L^n(\dot{G}_{PW}) = L(L^{n-1}(\dot{G}_{PW}))$  and  $\ddot{G}_{PW}$  is  $L(\ddot{G}_{PW})$ , and  $L^2(\ddot{G}_{PW}) = L(L(\ddot{G}_{PW}))$ , ...,  $L^n(\ddot{G}_{PW}) = L(L^{n-1}(\ddot{G}_{PW}))$

**Theorem (2.7):**  $S\bar{D}_{PW}, S\bar{D}_{PW}, S\dot{H}_{PW}$  and  $S\dot{H}_{PW}$  have network change symmetric Digraph-PW and

$$L(S\bar{D}_{PW}), L^2(S\bar{D}_{PW}) = L(L(S\bar{D}_{PW})), \dots, L^n(S\bar{D}_{PW}) = L(L^{n-1}(S\bar{D}_{PW})),$$

$$L(S\dot{H}_{PW}), L^2(S\dot{H}_{PW}) = L(L(S\dot{H}_{PW})), \dots, L^n(S\dot{H}_{PW}) = L(L^{n-1}(S\dot{H}_{PW})),$$

$$L(S\bar{D}_{PW}), L^2(S\bar{D}_{PW}) = L(L(S\bar{D}_{PW})), \dots, L^n(S\bar{D}_{PW}) = L(L^{n-1}(S\bar{D}_{PW})),$$

$$L(S\dot{H}_{PW}), L^2(S\dot{H}_{PW}) = L(L(S\dot{H}_{PW})), \dots, L^n(S\dot{H}_{PW}) = L(L^{n-1}(S\dot{H}_{PW})).$$

**Proof:** By the definition (2.7) the result follows.

**Theorem (2.8):** Every  $S\dot{D}_{PW}^r = (V, A_r^*, \Phi_r)$  there is  $L(S\dot{D}_{PW}^r), r = 1, 2, 3, \dots, \max d(v_i, v_j)$ .  
Moreover,  $L^2(S\dot{D}_{PW}^r) = L(L(S\dot{D}_{PW}^r)), \dots, L^n(S\dot{D}_{PW}^r) = L(L^{n-1}(S\dot{D}_{PW}^r))$ .

**Proof:** By theorem (1.20), by the definition (2.7) the result follows.

### 3. CONCLUSION

In this paper, we determined some new Networks. Furthermore, several theorems and results of these networks have been studied. In the future, we are interested in designing some new networks and then studying their topological indices which will be quite helpful in understanding their underlying topologies.

### 4. COMPLIANCE WITH ETHICAL STANDARDS

#### 4.1 AUTHORS' CONTRIBUTIONS

A. Alameri conceived of the presented idea, M. Alsharafi prepared the initial manuscript, W. A.M. Saeed developed the theory, A. Ghallab performed the computations, W. Yousef verified the analytical methods and A. Modabish supervised the findings of this work. All authors jointly worked on the results and they read and approved the final manuscript.

#### 4.2 AVAILABILITY OF DATA AND MATERIAL

No data were used to support this study.

#### 4.3 DECLARATION OF COMPETING INTEREST

The authors declare that they have no conflict of interest.

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## A Study on Network Graph-PW, Network Symmetric Digraph-PW, Change Network Graph-PW and Change Network Symmetric Digraph- PW

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